

# The meet-stalactic and meet-taiga monoids

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Stalactic monoids

$\text{ISt}, \text{rSt}$

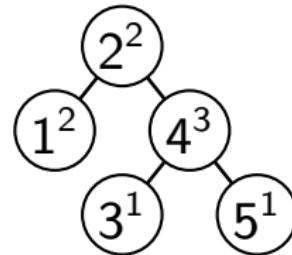
Stalactic tableaux

2	1	3	5	4
2	1	5		
1				

Taiga monoids

$\text{ITg}, \text{rTg}$

Binary search trees  
with multiplicities



# The left and right stalactic and taiga monoids

Let  $\text{ISt}$  and  $\text{rSt}$  denote the left and right stalactic monoids, and  $\text{ITg}$  and  $\text{rTg}$  denote the left and right taiga monoids.

## Definition

For  $u, v \in \mathbb{N}^*$ ,

$$u \equiv_{\text{ISt}} v \iff P_{\text{ISt}}(u) = P_{\text{ISt}}(v),$$

$$u \equiv_{\text{rSt}} v \iff P_{\text{rSt}}(u) = P_{\text{rSt}}(v),$$

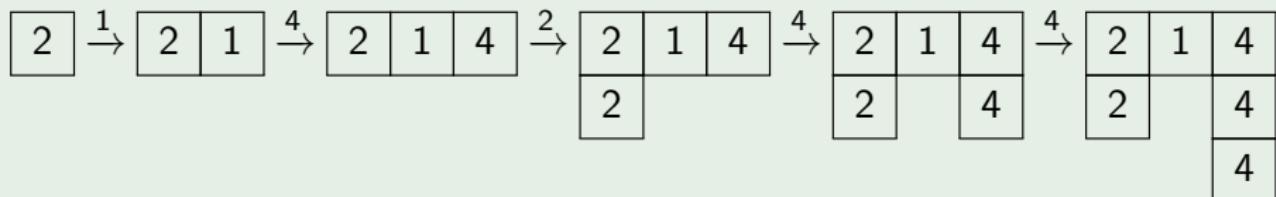
$$u \equiv_{\text{ITg}} v \iff P_{\text{ITg}}(u) = P_{\text{ITg}}(v), \text{ and}$$

$$u \equiv_{\text{rTg}} v \iff P_{\text{rTg}}(u) = P_{\text{rTg}}(v).$$

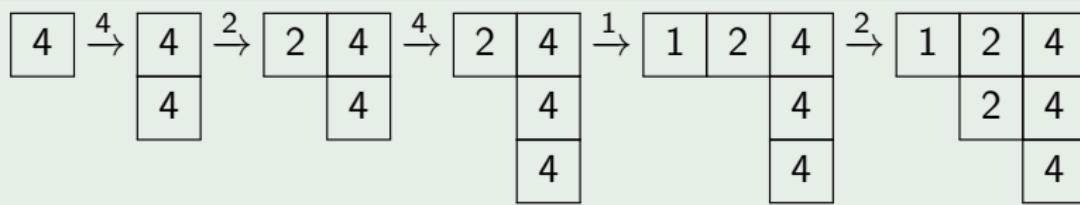
# The stalactic monoid

Given  $w = 214244 \in \mathbb{N}^*$ , we calculate  $P_{\text{St}}(w)$  and  $P_{\text{rSt}}(w)$ .

Example ( $P_{\text{St}}(214244)$ )



Example ( $P_{\text{rSt}}(214244)$ )



# Patience sorting tableau

## Definition

We define increasing [decreasing] patience sorting tableaux by example.  
Let,

$$A = \begin{array}{|c|c|c|} \hline 1 & 3 & 4 \\ \hline 2 & & 5 \\ \hline & 6 & \\ \hline \end{array}, \text{ and } B = \begin{array}{|c|c|c|} \hline 2 & 5 & 6 \\ \hline & 1 & 4 \\ \hline & & 3 \\ \hline \end{array}.$$

- ①  $A$  and  $B$  increase left-to-right in the first row.
- ②  $A$  increases top-to-bottom in each column, and
- ③  $B$  decreases top-to-bottom in each column.

So  $A$  is an *increasing* patience sorting tableau and  $B$  is a *decreasing* patience sorting tableau.

# Q-symbols for stalactic

Let  $w = 214244$ .

## Example

$$P_{\text{St}}(w) = \begin{array}{|c|c|c|} \hline 2 & 1 & 4 \\ \hline 2 & & 4 \\ \hline & 4 & \\ \hline \end{array} \text{ and } Q_{\text{St}}(w) = \begin{array}{|c|c|c|} \hline 1 & 2 & 3 \\ \hline 4 & & 5 \\ \hline & 6 & \\ \hline \end{array}.$$

## Example

$$P_{\text{rSt}}(w) = \begin{array}{|c|c|c|} \hline 1 & 2 & 4 \\ \hline 2 & 4 & \\ \hline & 4 & \\ \hline \end{array} \text{ and } Q_{\text{rSt}}(w) = \begin{array}{|c|c|c|} \hline 2 & 4 & 6 \\ \hline 1 & 5 & \\ \hline & 3 & \\ \hline \end{array}.$$

## Theorem

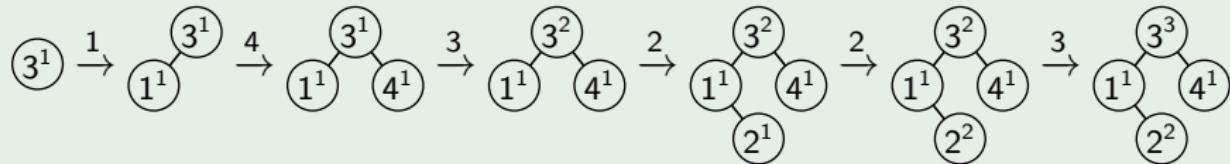
*The map  $w \mapsto (P_{\text{St}}(w), Q_{\text{St}}(w))$  [ $w \mapsto (P_{\text{rSt}}(w), Q_{\text{rSt}}(w))$ ] is a bijection between the elements of  $\mathbb{N}^*$  and the set formed by the pairs  $(T, S)$  where*

- ①  *$T$  is a stalactic tableau;*
- ②  *$S$  is an increasing [decreasing] patience-sorting tableau;*
- ③  *$T$  and  $S$  have the same shape.*

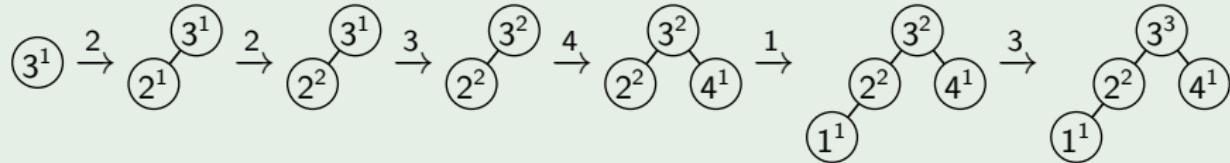
# The taiga monoid

Given  $w = 3143223 \in \mathbb{N}^*$ , we calculate  $P_{\text{ITg}}(w)$  and  $P_{\text{rTg}}(w)$ .

Example ( $P_{\text{ITg}}(3143223)$ )



Example ( $P_{\text{rTg}}(3143223)$ )

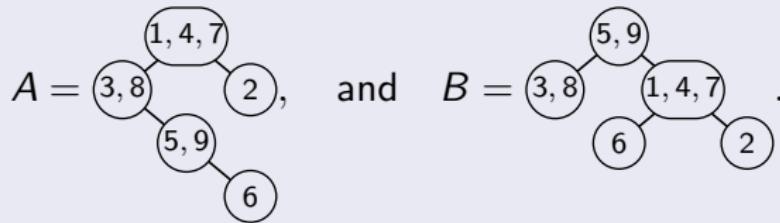


We call,  $P_{\text{ITg}}(w)$  and  $P_{\text{rTg}}(w)$ , *Binary search trees with multiplicities* (BSTM).

# Binary trees over sets

## Definition

We define an *increasing [decreasing] binary tree over sets* (BTS) by example. Let,



- ①  $A$  and  $B$  contain exactly the interval  $\{1, \dots, m\}$  for some  $m \in \mathbb{N}$ ,
- ② replacing each node in  $A$  with its minimum obtains an increasing tree,
- ③ replacing each node in  $B$  with its maximum obtains a decreasing tree.

So  $A$  is an increasing BTS and  $B$  is a decreasing BTS.

# Q-symbols for stalactic

Let  $w = 3143223$ .

## Example

$$P_{\text{ITg}}(w) = \begin{array}{c} 3^3 \\ \circlearrowleft \quad \circlearrowright \\ 1^1 \quad 4^1 \\ \circlearrowleft \quad \circlearrowright \\ 2^2 \end{array} \text{ and } Q_{\text{ITg}}(w) = \begin{array}{c} 1, 4, 7 \\ \circlearrowleft \quad \circlearrowright \\ 2 \quad 3 \\ \circlearrowleft \quad \circlearrowright \\ 5, 6 \end{array}$$

## Example

$$P_{\text{rTg}}(w) = \begin{array}{c} 3^3 \\ \circlearrowleft \quad \circlearrowright \\ 2^2 \quad 4^1 \\ \circlearrowleft \quad \circlearrowright \\ 1^1 \end{array} \text{ and } Q_{\text{rTg}}(w) = \begin{array}{c} 1, 4, 7 \\ \circlearrowleft \quad \circlearrowright \\ 5, 6 \quad 3 \\ \circlearrowleft \quad \circlearrowright \\ 2 \end{array}$$

## Theorem

*The map  $w \mapsto (P_{\text{ITg}}(w), Q_{\text{ITg}}(w))$  [ $w \mapsto (P_{\text{rTg}}(w), Q_{\text{rTg}}(w))$ ] a bijection between the elements of  $\mathbb{N}^*$  and the set formed by the pairs  $(T, S)$  where*

- ①  *$T$  is a BSTM;*
- ②  *$S$  is an increasing [decreasing] BTS such that the union of the sets labelling  $S$  is the interval  $[m]$ , where  $m$  is the sum of the multiplicities of  $T$ ;*
- ③  *$T$  and  $S$  have the same underlying binary tree shape;*
- ④ *the multiplicity of the  $i$ -th node of  $T$  is the cardinality of the set labelling the  $i$ -th node of  $S$ .*

# The meet-stalactic and meet-taiga monoid

Let  $mSt$  and  $mTg$  denote the meet-stalactic and meet taiga monoid.

## Definition

Let  $P_{mSt}(w) = (P_{lSt}(w), P_{rSt}(w))$ , and  $P_{mTg}(w) = (P_{lTg}(w), P_{rTg}(w))$ . Then, for  $u, v \in \mathbb{N}^*$ ,

$$u \equiv_{mSt} v \iff P_{mSt}(u) = P_{mSt}(v), \text{ and}$$

$$u \equiv_{mTg} v \iff P_{mTg}(u) = P_{mTg}(v).$$

## Example

We can equally define a  $Q$ -symbol for  $\text{mSt}$  by

$$Q_{\text{mSt}}(w) = (Q_{\text{ISSt}}(w), Q_{\text{rSt}}(w))$$

### Example

$$P_{\text{mSt}}(214244) = \left( \begin{array}{|c|c|c|} \hline 2 & 1 & 4 \\ \hline 2 & & 4 \\ \hline & 4 & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 1 & 2 & 4 \\ \hline & 2 & 4 \\ \hline & & 4 \\ \hline \end{array} \right) \text{ and}$$

$$Q_{\text{mSt}}(214244) = \left( \begin{array}{|c|c|c|} \hline 1 & 2 & 3 \\ \hline 4 & & 5 \\ \hline & 6 & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 2 & 4 & 6 \\ \hline & 1 & 5 \\ \hline & & 3 \\ \hline \end{array} \right)$$

# Twin stalactic tableaux

Let  $\text{cont}(T)$  be the multi-set containing every label in  $T$ . We say a column is *simple* if it only contains one block.

## Definition

Let  $T_L, T_R$  be stalactic tableaux. We say  $(T_L, T_R)$  are a *pair of twin stalactic tableaux* if

- ①  $\text{cont}(T_L) = \text{cont}(T_R)$ , and
- ② for each simple column labelled  $c$  and any column labelled  $d$ , if  $c$  is left of  $d$  in  $T_L$ , then  $c$  is left of  $d$  in  $T_R$ .

## Proposition

For any  $w \in \mathbb{N}^*$ , the meet-stalactic  $P$ -symbol  $(P_{\text{St}}(w), P_{\text{rSt}}(w))$  of  $w$  is a pair of twin stalactic tableaux.

# Twin patience-sorting tableaux

## Definition

Let  $(S_L, S_R)$  be a pair of (respectively) increasing and decreasing patience-sorting tableaux. We say  $(S_L, S_R)$  are a *pair of twin patience-sorting tableaux* if

- ① there exists a content preserving bijection  $\phi$  between the columns of  $S_L$  and the columns of  $S_R$ ;
- ② for each simple column  $c$  and any column  $d$  in  $S_L$ , if  $c$  appears to the left of  $d$  in  $S_L$  then  $\phi(c)$  appears to the left of  $\phi(d)$  in  $S_R$ .

## Proposition

For any  $w \in \mathbb{N}^*$ , the meet-stalactic  $Q$ -symbol  $(Q_{\text{St}}(w), Q_{\text{rSt}}(w))$  of  $w$  is a pair of twin patience-sorting tableaux.

## Theorem

The map  $w \mapsto (P_{\text{mSt}}(w), Q_{\text{mSt}}(w))$  is a bijection between the elements of  $\mathbb{N}^*$  and the set formed by the pairs  $((T_L, T_R), (S_L, S_R))$  where

- ①  $(T_L, T_R)$  is a pair of twin stalactic tableaux;
- ②  $(S_L, S_R)$  is a pair of twin patience-sorting tableaux with bijection  $\phi$ ;
- ③  $(T_L, T_R)$  and  $(S_L, S_R)$  have the same shape;
- ④ For each column  $c$  in  $S_L$ , the column in  $T_L$  in the same position as  $c$  and the column in  $T_R$  in the same position as  $\phi(c)$ , have the same content.

# Meet-taiga example

We can equally define a  $Q$ -symbol for  $mTg$  by

$$Q_{mTg}(w) = (P_{lTg}(w), P_{rTg}(w))$$

## Example

$$P_{mTg}(3143223) = \left( \begin{array}{c} \text{Diagram 1: } \begin{array}{c} 3^3 \\ \diagup \quad \diagdown \\ 1^1 \quad 4^1 \\ \diagdown \quad \diagup \\ 2^2 \end{array}, \quad \text{Diagram 2: } \begin{array}{c} 3^3 \\ \diagup \quad \diagdown \\ 1^1 \quad 4^1 \\ \diagdown \quad \diagup \\ 2^2 \end{array} \end{array} \right) \text{ and}$$
$$Q_{mTg}(3143223) = \left( \begin{array}{c} \text{Diagram 1: } \begin{array}{c} 1, 4, 7 \\ \diagup \quad \diagdown \\ 2 \quad 3 \\ \diagdown \quad \diagup \\ 5, 6 \end{array}, \quad \text{Diagram 2: } \begin{array}{c} 1, 4, 7 \\ \diagup \quad \diagdown \\ 5, 6 \quad 3 \\ \diagdown \quad \diagup \\ 2 \end{array} \end{array} \right)$$

# Twin binary search trees with multiplicity

## Definition

Let  $T_L, T_R$  be BTSMs. We say  $(T_L, T_R)$  is a *pair of twin binary trees with multiplicities* (pair of twin BTMs), if for all  $i$ ,

- ①  $\text{cont}(T_L) = \text{cont}(T_R)$ .
- ② if the  $i$ -th node of  $T_L$  has multiplicity 1 and has a left (resp. right) child then the  $i$ -th node of  $T_R$  does not have a left (resp. right) child.

## Proposition

For any  $w \in \mathbb{N}^*$ , the meet-taiga  $P$ -symbol  $(P_{\text{ITg}}(w), P_{\text{rTg}}(w))$  of  $w$  is a pair of twin BTMs.

# Twin binary trees over sets

## Definition

Let  $(S_L, S_R)$  be a pair of (respectively) increasing and decreasing BTSSs. We say  $(S_L, S_R)$  are a *pair of twin binary trees over sets* (pair of twin BTSSs) if, for all  $i$ ,

- ① the  $i$ -th node of  $S_L$  has the same label as the  $i$ -th node of  $S_R$ ;
- ② if the  $i$ -th node of  $S_L$  is labelled by a set of cardinality 1 and has a left (resp. right) child, then the  $i$ -th node of  $S_R$  does not have a left (resp. right) child.

## Proposition

For any  $w \in \mathbb{N}^*$ , the meet-taiga  $Q$ -symbol  $(Q_{\text{LTg}}(w), Q_{\text{RTg}}(w))$  of  $w$  is a pair of twin BTSSs.

## Theorem

*The map  $w \mapsto (P_{\text{mTg}}(w), Q_{\text{mTg}}(w))$  is a bijection between the elements of  $\mathbb{N}^*$  and the set formed by the pairs  $((T_L, T_R), (S_L, S_R))$  where*

- ①  *$(T_L, T_R)$  is a pair of twin BSTMs;*
- ②  *$(S_L, S_R)$  is a pair of twin BTSs;*
- ③  *$(T_L, T_R)$  and  $(S_L, S_R)$  have the same underlying pair of binary trees shape;*
- ④ *the multiplicity of the  $i$ -th node of  $T_L$  (resp.  $T_R$ ) is the cardinality of the set labelling the  $i$ -th node of  $S_L$  (resp.  $S_R$ ).*

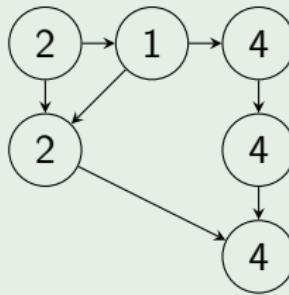
# Size of stalactic classes

## Example

Consider the meet-stalactic tableau,

$$T = P_{\text{mSt}}(214244) = \left( \begin{array}{|c|c|c|} \hline 2 & 1 & 4 \\ \hline 2 & & 4 \\ \hline & 4 & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 1 & 2 & 4 \\ \hline & 2 & 4 \\ \hline & & 4 \\ \hline \end{array} \right)$$

Then, we define the following poset:  $\psi_{\text{mSt}}(T) =$



## Proposition

Let  $w \in \mathbb{N}^*$ . Then, each word in  $[w]_{\text{mSt}}$  is in bijection with a linear extension of  $\psi_{\text{mSt}}(P_{\text{mSt}}(w))$ .

## Proposition

There exists an algorithm to compute the number of linear extensions of  $\psi_{\text{mSt}}(T)$  has time complexity  $\mathcal{O}(n^{2k-2}k!)$  where  $n$  is the number of nodes and  $k$  is the size of the support.

## Theorem

Let  $k \geq 2$ ,  $V = (v_1, \dots, v_k) \in \mathbb{N}_0^k$ ,  $B = (b_1, \dots, b_{k-1}) \in \mathbb{N}_0^{k-1}$  and  $\sigma$  be a permutation of  $[k]$ . Then, when  $\sigma_1 < \sigma_2$ ,  $\mathcal{L}(G[V; B; \sigma])$  is equal to

$$\sum_{\substack{M \in \mathbb{N}_0^{\sigma_2 - \sigma_1} \\ 0 \leq \|M\|_1 \leq v_{\sigma_1}}} \mathcal{L}_M \cdot \binom{v_{\sigma_2} - 1 + v_{\sigma_1} - \|M\|_1}{v_{\sigma_1} - \|M\|_1} \prod_{i=1}^{\sigma_2 - \sigma_1} \binom{b_{\sigma_1 - 1 + i} + m_i}{m_i},$$

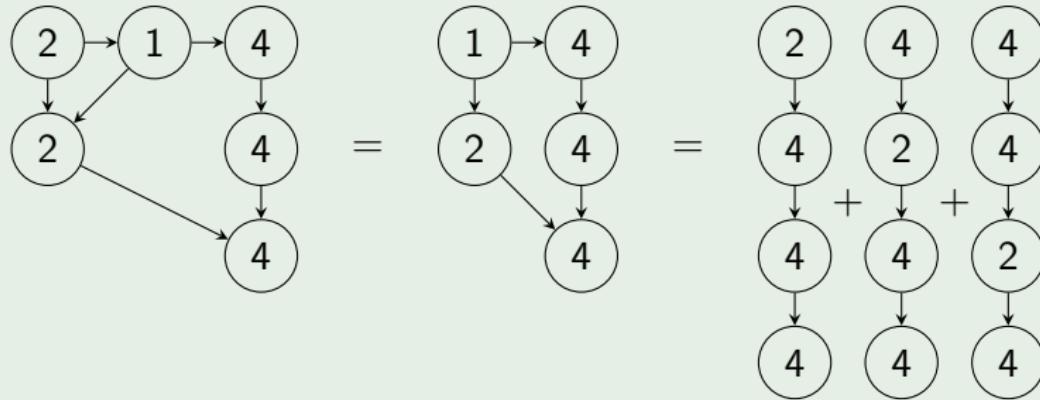
where  $\mathcal{L}_M = 1$  if  $k = 2$  and  $\mathcal{L}_M = \mathcal{L}(G[f_{k, \sigma_1, \sigma_2, M}(V, B); \text{Std}(\sigma_2 \sigma_3 \cdots \sigma_k)])$  otherwise, and, when  $\sigma_2 < \sigma_1$ , is equal to

$$\sum_{\substack{M \in \mathbb{N}_0^{\sigma_1 - \sigma_2} \\ 0 \leq \|M\|_1 \leq v_{\sigma_2} - 1}} \mathcal{L}'_M \cdot \binom{v_{\sigma_1} + v_{\sigma_2} - 1 - \|M\|_1}{v_{\sigma_2} - 1 - \|M\|_1} \prod_{i=1}^{\sigma_1 - \sigma_2} \binom{b_{\sigma_2 - 1 + i} + m_i}{m_i},$$

where  $\mathcal{L}'_M = 1$  if  $k = 2$  and  $\mathcal{L}'_M = \mathcal{L}(G[f_{k, \sigma_2, \sigma_1, M}(V, B); \text{Std}(\sigma_1 \sigma_3 \cdots \sigma_k)])$  otherwise.

# The Algorithm

## Example



So,  $[214244]_{\text{mSt}} = 3$ .

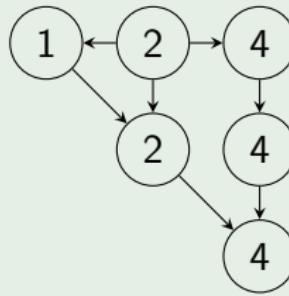
# Size of taiga classes

## Example

Consider the meet-taiga tableau,

$$T = P_{\text{mTg}}(214244) = \left( \begin{array}{c} \text{1}^1 \text{2}^2 \text{4}^3, \\ \text{2}^2 \text{4}^3 \end{array} \right)$$

Then, we define the following poset:  $\psi_{\text{mTg}}(T) =$



# Size of taiga classes

## Proposition

Let  $w \in \mathbb{N}^*$ . Then, each word in  $[w]_{\text{mTg}}$  is in bijection with a linear extension of  $\psi_{\text{mTg}}(P_{\text{mTg}}(w))$ .

## Theorem

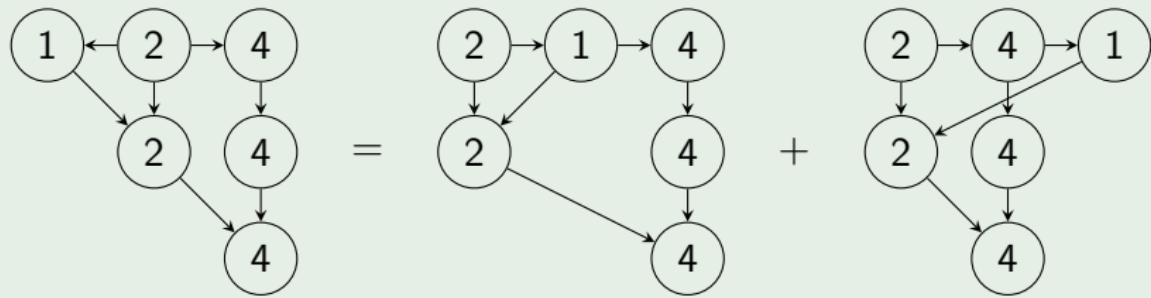
Let  $Q_{\text{mTg}}(w) = T = (T_L, T_R)$ . Then,

$$\mathcal{L}(\psi_{\text{mTg}}(T)) = \sum_L \sum_R \mathcal{L}(\text{mSt}(P_{\text{mSt}}(w_{L,R})))$$

where the first sum is over all linear extensions  $L$  of  $\Delta(T_L)$  and the second sum is over all linear extensions  $R$  of  $\nabla(T_R, p_L)$  where  $a < x$  in  $p_L$  if and only if  $a$  is simple in  $(T_L, T_R)$  and  $a < x$  in  $L$ .

# The Algorithm

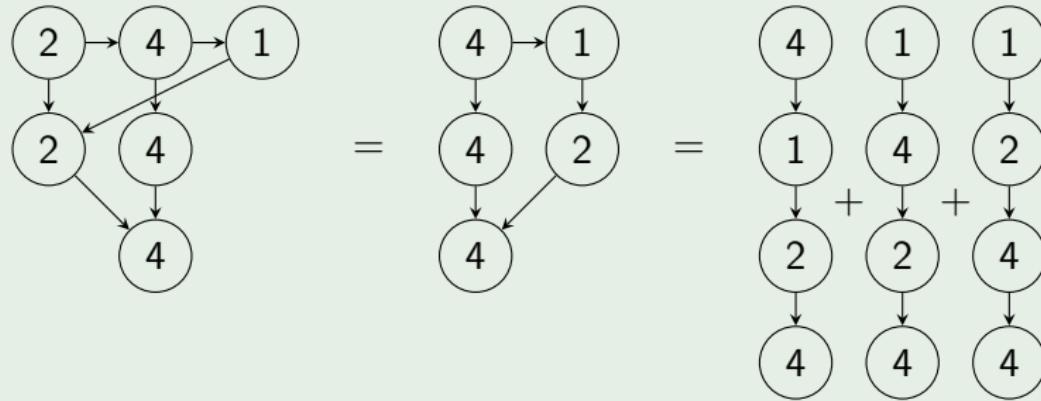
## Example



$$\text{So, } [214244]_{mTg} = [214244]_{mSt} + [241244]_{mSt}.$$

# The Algorithm

## Example



So,  $[214244]_{mTg} = 3 + 3 = 6$ .

## Proposition

*There exists an algorithm to compute the number of linear extensions of  $\psi_{mTg}(T)$  has time complexity  $\mathcal{O}(n^{2k-2}(k!)^3)$  where  $n$  is the number of nodes and  $k$  is the size of the support.*